

# Example Applications of Finite State Machines

Data structures and algorithms  
for Computational Linguistics III

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Seminar für Sprachwissenschaft

Winter Semester 2018–2019

# Applications of finite-state methods

- Finite state methods are attractive for formal and computational reasons
- They are applied in a vast diversity of fields
  - Electronic circuit design
  - Workflow management
  - Games
  - Pattern matching
  - Tokenization, stemming
  - Morphological analysis
  - Chunking
  - ...
- This lecture
  - FSA for pattern matching
  - FSA for storing a lexicon
  - Finite-state morphology

# Finite state automata

## a refresher

- An FSA recognizes and generates a regular language, also equivalent to regular expressions
- FSA are closed under
  - Concatenation
  - Kleene star
  - Union
  - Intersection
  - Complement
  - Reversal
- Two types:
  - DFA single transition from each state on each input symbol
  - NFA transitions to possibly multiple states on a single input symbol, or without consuming an input symbol ( $\epsilon$ -NFA)
- Every FSA has a unique minimal DFA
  - For every NFA there is a DFA that accepts the same regular language (determinization)
  - A DFA can be minimized to equivalent DFA with minimum nodes (minimization)

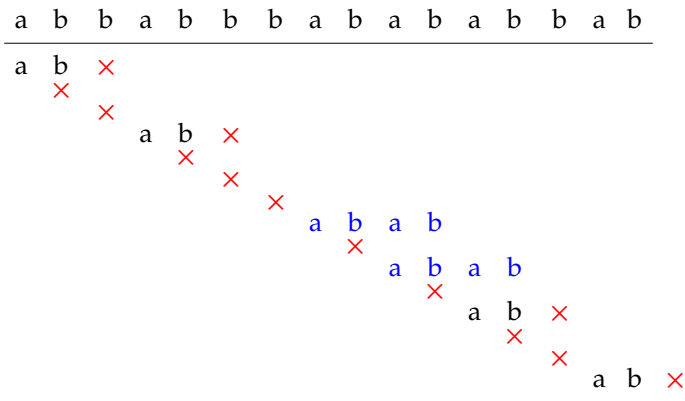
# Finite state transducers

## a refresher

- FST transitions are defined on a pair of input–output symbols
- An FST moves between the states on the input symbol, while outputting the output symbol
- FSTs define a regular relation
- FSTs are closed under
  - Concatenation
  - Union
  - Inversion
  - Kleene star
  - Reversal
  - Composition
- Not all FSTs can be determinized

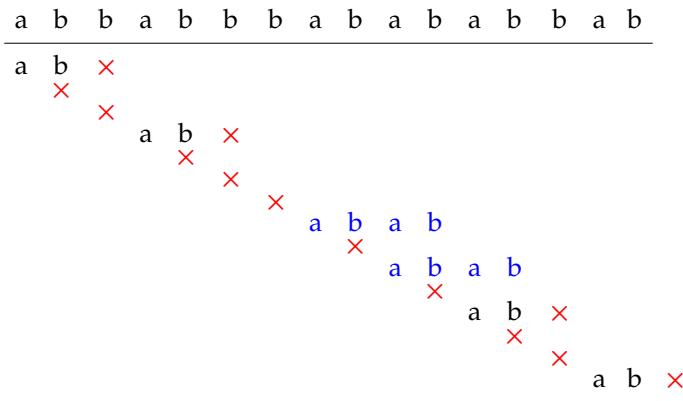
# Naive string match

Example: searching 'abab' in 'abbabbbabababbab'



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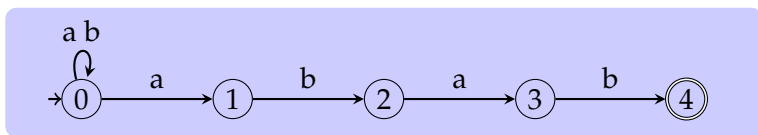


Note the wasted effort after a partial match.

# String matching with an NFA

## Another solution

Consider running the following NFA over the string.

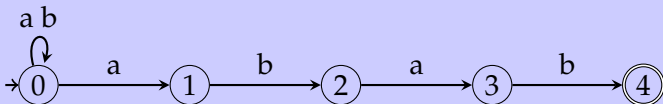


- The NFA will be in the accepting state when last four letters processed matches abab (including overlapping matches)

# String matching with an NFA

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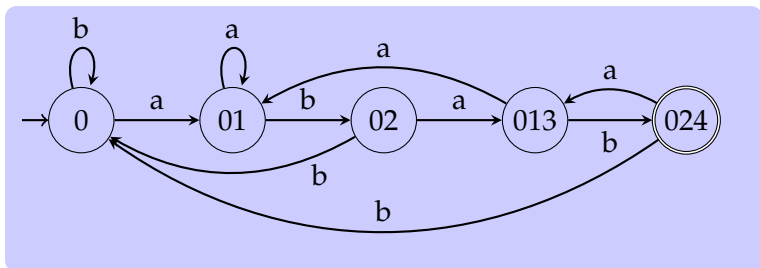


- The NFA will be in the accepting state when last four letters processed matches abab (including overlapping matches)
- Is this faster than the naive algorithm?



# DFA version

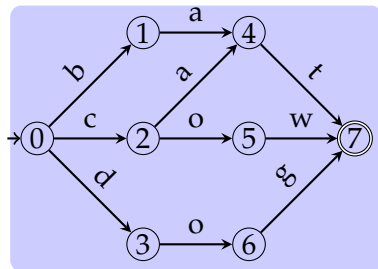
## Knuth-Morris-Pratt (KMP) algorithm



- DFA processes every input symbol only once
- The resulting DFA has the same number of states (generally, not much larger than the NFA)
- Approach generalizes to arbitrary regular expressions without additional computational cost

## Finite state lexicons

- FSA are an efficient way to store lexicons
- One can start from NFA for individual words, and minimize/determinize the union of them
- Or there are algorithms for constructing finite-state lexicons incrementally



# Morphology

## some definitions

**Morpheme** is an abstract linguistic unit, often defined as smallest meaningful or grammatical unit.

Morphemes make up words

**Root** of a word is a *free* morpheme, often carrying the semantic information

**Derivational** morphemes change the meaning of a word, sometimes changing the POS

**Inflectional** morphemes change the syntactic properties of words

**Lemma** of a word is its 'citation' form, what you look up in a lexicon

**Stem** of a (possibly derived) word is the common string shared by all morphologically related forms

# Morphological typology

Languages of the world behave differently with respect to how words are formed.

- *Isolating* languages have little or no morphology, all words are simple (e.g., Vietnamese, Chinese)
- *Analytic* languages have little or no inflectional morphology (e.g., English)
- *Synthetic* languages have rich morphological system
  - In *agglutinative* languages each morpheme has a single function (e.g., Finnish, Turkish)
  - In *inflecting/fusional* a single morpheme indicates multiple functions (e.g., Latin, Russian)
  - *Polysynthetic* languages may pack multiple ‘words’ in a single word (e.g., Ainu, Chukchi)

Note that these are tendencies.

# Where do morphemes go

- Affixation:  
*attach* → *un-attach-ed*
- Infixes:  
*aussteigen* → *auszusteigen*
- Circumfixation:  
*spiel* → *gespielt*
- Root-pattern morphology:  
*ktb* → *kitāb* 'book'  
*ktb* → *kātib* 'writer'  
(Arabic)
- Reduplication:  
*orang* 'person' → *orang-orang* 'people'  
(some Austronesian languages)

# Interaction of morphology and phonology

## or morphology and orthography

Morphology and phonology/orthography interact. A few examples:

- *dog-s*, but *fox-es*
- *city* → *citi-es*
- *stop* → *stopping*
- *panic* → *panick-ed*
- *goose* → *geese*
- Vowel harmony
 

<i>ev</i> ‘house’	→	<i>ev-ler</i> ‘houses’	(Turkish)
<i>oda</i> ‘room’	→	<i>oda-lar</i> ‘rooms’	

## Two-level morphology

- We assume that there are two ‘levels’ of representation
  - A *surface* representation which is what we hear or see
  - An *underlying*, an abstract representation for the word

Surface:	cat	s
Underlying:	cat	⟨PL⟩

- An FST is used to map the underlying representation to the surface representation (generation)
- If we run the FST in the inverse direction, we get an analysis
- Often the FST is a complex combination of many small FSA or FSTs

# Two-level morphology

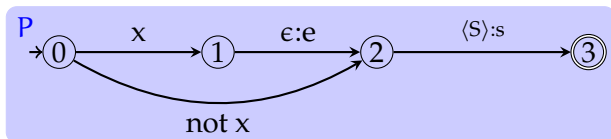
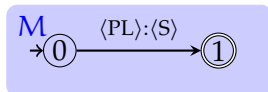
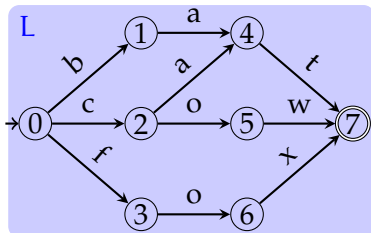
## a typical architecture

- Typically, lexicon is converted to FSA
- Concatenated (or composed) with morphological rules (affixation, applying templates, ...)
- The result is composed with phonological/orthographic alternations
- The phonological/orthographic rules can be designed as cascades (composition), or can be applied in parallel



# Two-level morphology

a (simplified) example



Generator:  $LM \circ P$

Analyzer:  $(LM \circ P)^{-1}$

## How to specify morphological analyzers

- Lexicons are easiest to specify as lists of (root) words
  - cat
  - dog
  - fox
  - ...
- For affixation, regular expressions (or regular rewrite rules)
  - $N_{\text{plu}} \rightarrow N \langle \text{PL} \rangle : \langle \text{S} \rangle$
- For phonological/orthographic alternations context sensitive rules
  - $\langle \text{S} \rangle \rightarrow \text{es} / \text{x}_-$
- There are a few standard languages for specifying morphological analyzers
  - SFST
  - Xerox languages: XFST, Twolc, lexc
  - OpenFST OpenGRM (more general purpose)

# XFST

A quick reference some common notation/operations

?	any symbol
$\emptyset$	empty string ( $\epsilon$ )
(a)	optional a
[a b]	grouping
a*	Kleene star
a+	Kleene plus
a b	concatenation
a&b	intersection
a b	union
$\sim$ b	complement
a-b	difference
{cat}	concatenation of c a t
a:b	FST rule with input 'a' and output 'b'
a .o. b	compose a with b
a -> b	unconditionally replace a to b

## XFST (cont.)

A quick reference some common notation/operations

- a (->)b                      optionally replace a to b
- a -> b || c \_              replace a to b only after c
- a -> b || c \_ d            replace a to b only after c and before d

- There are (at least) two free implementations of xfst
  - Foma
  - hfst-xfst (part of HFST)
- You will receive a separate ‘tutorial’ (and an exercise) on working with xfst and lexc

## Tools of the trade

Some of the practical, feely-available, tools (with an emphasis on ones targeted for CL) include:

- Gertjan van Noord's **FSA tools**
- **OpenFST**: a general purpose finite state library
- **Helsinki finite-state technology** (HFST): library tools from University of Helsinki
- **Foma**: a re-implementation of Xerox's xfst, a language/toolbox for defining/manipulating FST
- **SFST** another language/toolbox for defining/manipulating FSTs

## Wrapping up

- Finite-state tools are commonly used in a number of CL task
- There are off-the-shelf free tools

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Next:

- Dependency grammars and dependency parsing
- Constituency parsing

## References / additional reading material

- Jurafsky and Martin (2009, Ch. 3)
- Roche and Schabes (1997) includes more examples of FSTs used for NLP
- The Xerox languages and tools are described in Beesley and Karttunen (2003)
- **HFST** and **Foma** web pages include some documentation and (links to) tutorials



## References / additional reading material (cont.)

- Beesley, Kenneth R. and Lauri Karttunen (2003). "Finite-state morphology: Xerox tools and techniques". In: *CSLI, Stanford*.
- Hulden, Mans (2009). "Foma: a finite-state compiler and library". In: *Proceedings of the 12th Conference of the European Chapter of the Association for Computational Linguistics*. Association for Computational Linguistics, pp. 29–32.
- Jurafsky, Daniel and James H. Martin (2009). *Speech and Language Processing: An Introduction to Natural Language Processing, Computational Linguistics, and Speech Recognition*. second. Pearson Prentice Hall. ISBN: 978-0-13-504196-3.
- Lindén, Krister, Erik Axelsson, Senka Drobac, Sam Hardwick, Juha Kuokkala, Jyrki Niemi, Tommi A. Pirinen, and Miikka Silfverberg (2013). "HFST — A System for Creating NLP Tools". In: *Systems and Frameworks for Computational Morphology*. Ed. by Cerstin Mahlow and Michael Piotrowski. Berlin, Heidelberg: Springer Berlin Heidelberg, pp. 53–71.
- Roche, Emmanuel and Yves Schabes (1997). *Finite-state Language Processing*. A Bradford book. MIT Press. ISBN: 9780262181822.